



On the Role of Coherence in Shor's algorithm

Introduction to QRTs and applications

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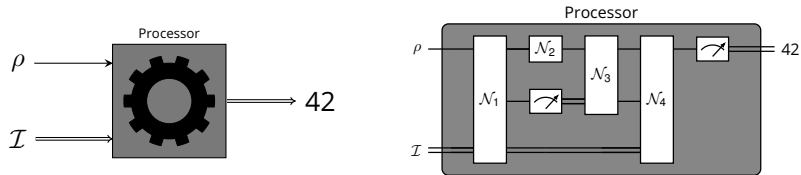
QLOC Seminar
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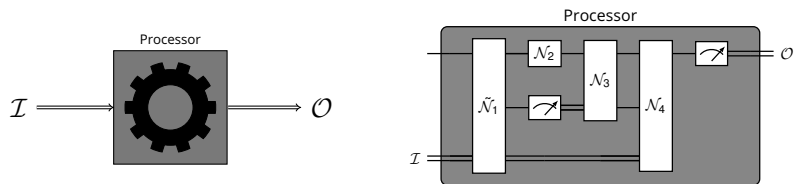
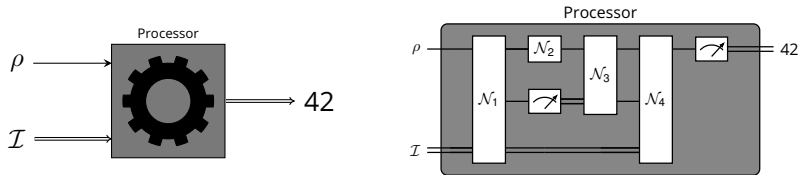
In this episode...

- 1 Quantum resource theories
- 2 Coherence as a resource in quantum algorithms

Motivation



Motivation



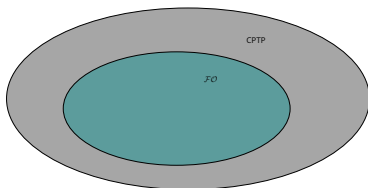
What fuels the processor?

Operations and free operations

Definition

Let $\mathcal{FO}$ be a the set of maps which assigns a map $\mathcal{FO}(A \rightarrow B) \subset \text{CPTP}(A \rightarrow B)$ to any two systems A and B . The set $\mathcal{FO}$ is referred to as the set of *free operations* if

- 1 For any system A , the set $\mathcal{FO}(A \rightarrow A)$ contains the identity channel id_A .
- 2 For any systems A, B, C , if $\Phi_1 \in \mathcal{FO}(A \rightarrow B)$ and $\Phi_2 \in \mathcal{FO}(B \rightarrow C)$ the composition of maps is closed, i.e., $\Phi_1 \circ \Phi_2 \in \mathcal{FO}(A \rightarrow C)$.



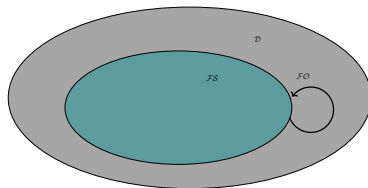
Free states & the golden rule

Definition

The set of *free states* is defined as $\mathcal{F} = \mathcal{FO}(\mathbb{C} \rightarrow A)$ for any system A .

Golden rule

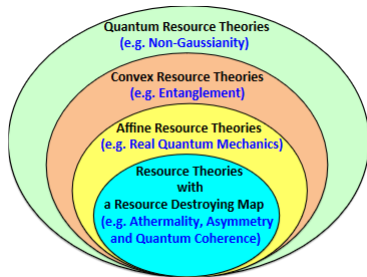
For any A and B , if
 $\Phi \in \mathcal{FO}(A \rightarrow B)$ and $\rho \in \mathcal{F}$,
then $\Phi(\rho) \in \mathcal{F}$.



Definition

The tuple $\mathcal{R} = (\mathcal{F}, \mathcal{FO})$ defines a *static* resource theory.

QRT hierarchy & examples



Borrowed from Chitambar and Gour

- Entanglement

$$\Lambda_{\text{free}}(\cdot) = \sum_k \left(\bigotimes_i M_{k,i}^{A_i} \right) (\cdot) \left(\bigotimes_i M_{k,i}^{A_i} \right)^\dagger$$

$$\rho \in \text{SEP} \Leftrightarrow \rho = \sum_k p_k \bigotimes_i \rho_{i,k}^{A_i}$$

- Imaginary QM

$$\Lambda \in \mathcal{FO} \Leftrightarrow J_\Lambda = (\text{id} \otimes \Lambda)|\Omega\rangle\langle\Omega| \quad \text{real}$$

$$\rho \in \mathcal{RS} \Leftrightarrow \rho \quad \text{real}$$

- Coherence: Resource destroying map Δ

$$\Delta(\cdot) = \sum_i |i\rangle\langle i|(\cdot)|i\rangle\langle i| \quad \Lambda \in \mathcal{FO} \Leftrightarrow \Lambda \circ \Delta = \Delta \circ \Lambda \circ \Delta \quad \rho \in \mathcal{I} \Leftrightarrow \Delta(\sigma) = \sigma$$

Chitambar, Gour (2018). Rev. Mod. Phys. 91, 025001
 Horodecki family (2009). Rev. Mod. Phys. 81, 865
 Hickey, Gour (2018) J. Phys. A: Math. Theor. 51 414009

Liu, Hu, Lloyd (2017). Phys. Rev. Lett. 118, 060502
 Streltsov, Adesso, Plenio (2017). Rev. Mod. Phys. 89, 041003

Super-operators and free super-operators

Lemma

A linear map $S : CPTP(A \rightarrow B) \rightarrow CPTP(C \rightarrow D)$ can be represented as

$$S[\Lambda] = \Theta \Leftrightarrow \Theta_2(\Lambda \otimes \mathbb{1})\Theta_1 = \Theta.$$

This map is referred to as a super-operation.

- Sequential and parallel concatenation is free

$$\Phi_2\Phi_1 \in \mathcal{F} \quad \text{and} \quad \Phi_1 \otimes \Phi_2 \in \mathcal{F} \quad \forall \Phi_1, \Phi_2 \in \mathcal{F}$$

Definition

A super-operation is called *free* if any free operation is transformed into a free operations, i.e., S free $\Leftrightarrow S[\Lambda] = \Phi_2(\Lambda \otimes \mathbb{1})\Phi_1 \quad \Phi_1, \Phi_2 \in \mathcal{F}$.

Pre-order and resource measures

$$\begin{array}{ll} \rho_2 \preceq \rho_1 \Leftrightarrow \exists \Phi \in \mathcal{FO} : \Phi(\rho_1) = \rho_2 & \Theta_2 \preceq \Theta_1 \Leftrightarrow \exists S \in \mathcal{FSO} : S[\Theta_1] = \Theta_2 \\ \rightarrow \rho_1 \text{ more valuable} & \rightarrow \Theta_1 \text{ more valuable} \end{array}$$

Pre-order and resource measures

$$\rho_2 \preceq \rho_1 \Leftrightarrow \exists \Phi \in \mathcal{FO} : \Phi(\rho_1) = \rho_2 \quad \Theta_2 \preceq \Theta_1 \Leftrightarrow \exists S \in \mathcal{FSO} : S[\Theta_1] = \Theta_2$$

$\rightarrow \rho_1$ more valuable

$\rightarrow \Theta_1$ more valuable

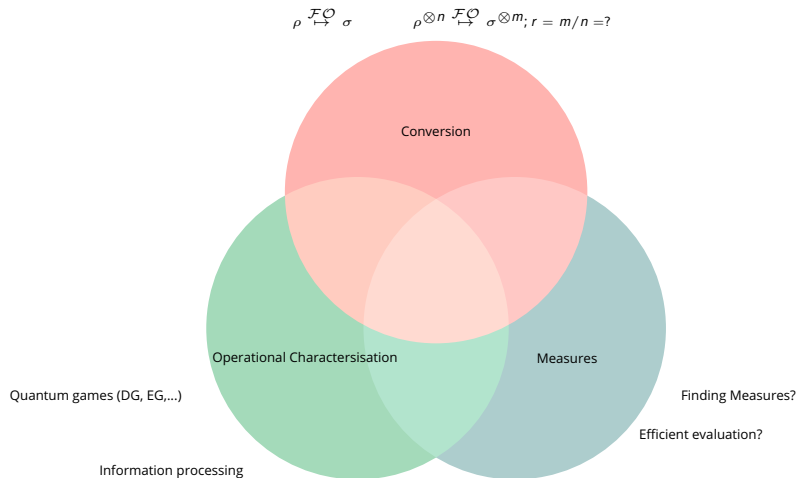
Does a Φ or S exist?
 $\rightarrow$ Resource measures

Definition

A positive functional M is called a static (dynamic) resource measure if it satisfies:

- 1 Monotonicity: $M(\Phi(\rho)) \leq M(\rho) \quad \forall \rho \quad \left(M(S[\Theta]) \leq M(\Theta) \quad \forall \Theta \right)$
- 2 Faithfulness: $M(\rho) = 0 \Leftrightarrow \rho \in \mathcal{F} \quad \left(M(\Theta) = 0 \Leftrightarrow \Theta \in \mathcal{FO} \right)$
- 3 Convexity: M is convex (for convex QRT)

Problems to study



Chitambar, Gour (2018). Rev. Mod. Phys. 91, 025001

Streltsov, Adesso, Plenio (2017) Rev. Mod. Phys. 89,041003

Concrete example: Coherence

$$\{|0\rangle, |1\rangle\} \quad \text{and} \quad \begin{aligned} |+\rangle &= |0\rangle + |1\rangle, \\ |-\rangle &= |0\rangle - |1\rangle \end{aligned}$$

Concrete example: Coherence

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- 1 Fix a preferred basis $\{|i\rangle\}_i$ and

$$\Delta(\cdot) = \sum_i |i\rangle\langle i|(\cdot)|i\rangle\langle i|$$

- 2 RNG set of free operations $\mathcal{MIO}$

$$\Phi \in \mathcal{MIO} \Leftrightarrow \Phi\Delta = \Delta\Phi\Delta$$

Application: Playing discrimination games

Discrimination game

Let $\{p_i, \mathcal{N}_i\}_{i=1\dots N}$ be a black-box operation that implements a channel $\mathcal{N}_i$ with probability p_i . The probability to successfully discriminate which channel was applied is given by

$$P^{\text{succ}}(\rho) = \max_{\{M_n\}} \sum_n p_n \text{Tr} [\mathcal{N}_n(\rho) M_n]$$

Application: Playing discrimination games

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$$P^{\text{succ}}(\rho) = \max_{\{M_n\}} \sum_n p_n \text{Tr} [\mathcal{N}_n(\rho) M_n]$$

Lemma

For $\{\frac{1}{N}, \mathcal{U}_n\}$, with $U_n = \sum_m e^{i\phi_n m} |m\rangle\langle m|$, any non-free state provides an advantage, i.e.,

$$P^{\text{succ}}(\rho) \geq P^{\text{succ}}(\sigma) \quad \forall \sigma \in \mathcal{F}.$$

Proof.

Take $\rho = V\sigma V^\dagger$ and since $P^{\text{succ}}(\sigma_1) = P^{\text{succ}}(\sigma_2) \forall \sigma_1, \sigma_2$

$$\begin{aligned} P^{\text{succ}}(\rho) &= \max_{\{M_n\}} \sum_n p_n \text{Tr} \left[UV\sigma V^\dagger U^\dagger M_n \right] \\ &= \max_{\{M_n\}} \sum_n p_n \text{Tr} \left[\sigma(V^\dagger U^\dagger M_n UV) \right] \\ &\geq P^{\text{succ}}(\sigma) \end{aligned}$$



Proof.

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□

Proposition

The optimal advantage for any discrimination game $\{p_n, \mathcal{U}_n\}$ is quantified by

$$\max_{\{p_n\}} \frac{P^{\text{succ}}(\rho)}{P^{\text{succ}}(\sigma \in \mathcal{I})} = 1 + C(\rho) \quad \text{where} \quad C(\rho) = \min_{\tau \in \mathcal{D}} \left\{ r \geq 0 : \frac{\rho + r\tau}{1+r} =: \sigma \in \mathcal{I} \right\}.$$

Motive to name these frameworks as resource theories.

Why resource theories are important?

Coherence as a Resource for Shor's Algorithm

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⁴*IFLP-CONICET, Departamento de Física, Facultad de Ciencias Exactas, Universidad Nacional de La Plata, C.C. 67, La Plata 1900, Argentina*



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Dario



Mauricio



Martin

Dynamical QRT of Coherence I

- Resource destroying map

$$\Delta(\rho) = \sum_i |i\rangle\langle i| \rho |i\rangle\langle i|$$

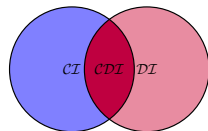
- Free Operations

① Creation-incoherent:

$$\Phi\Delta = \Delta\Phi\Delta$$

② Detection-incoherent:

$$\Delta\Phi = \Delta\Phi\Delta$$



Streltsov, Adesso, Plenio (2017). Rev. Mod. Phys. 89, 041003

Theurer, Egloff, Plenio (2019). Phys. Rev. Lett. 122, 190405

Saxena, Chitambar, Gour (2020). Phys. Rev. Research 2, 023298

Dynamical QRT of Coherence II

- Incoherent measurements $M = \{P_n\}$

$$\text{Tr}[P_n \rho] = \text{Tr}[P_n \Delta \rho] \quad \forall n, \rho$$

$$M \text{ free} \Leftrightarrow P_n = \sum_i P_i^{(n)} |i\rangle\langle i|$$

- Incoherent states: $\Delta(\sigma) = \sigma$

Dynamical QRT of Coherence II

- Incoherent measurements $M = \{P_n\}$

$$\text{Tr}[P_n \rho] = \text{Tr}[P_n \Delta \rho] \quad \forall n, \rho$$

$$M \text{ free} \Leftrightarrow P_n = \sum_i P_i^{(n)} |i\rangle\langle i|$$

- Incoherent states: $\Delta(\sigma) = \sigma$

Creation-Incoherent

- 1 MIO free operations
- 2 $MIOS$ free super-operations
- 3 $\mathcal{I}$ inc. states

Detection-Incoherent

- 1 DI free operations
- 2 DIS free super-operations
- 3 $\mathcal{IM}$ inc. measurements

Previous works?

Shor's coup

Goal: Factor integer N

⇒ Solve order-finding via $\{\frac{1}{r}, \mathcal{E}_j\}_{j=0..r-1}$

$$\mathcal{E}_j \equiv \sum_n e^{2\pi i \frac{j}{r}} |n\rangle\langle n|$$

Recover integer r ?

⇒ Efficient post-processing via CFA

Shor's coup

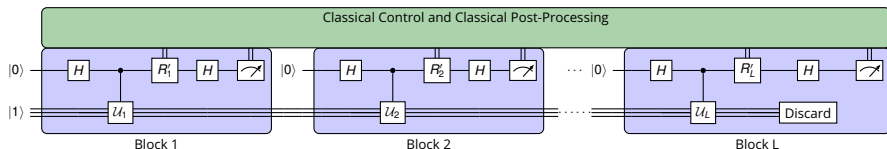
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Recover integer r ?

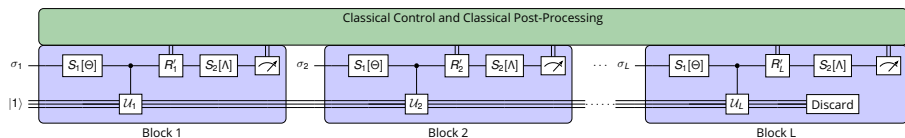
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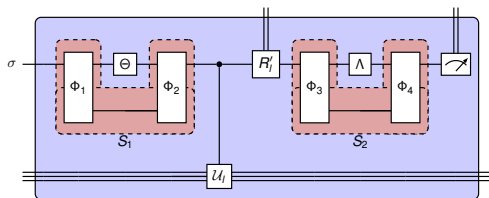
$$N^2 < 2^L < 2N^2$$

A partially coherent protocol

- Keep all incoherent and classical operations
⇒ Replace all others



Success Probability



$$P^{\text{succ}}(\Theta, \Lambda) = \max_{\sigma \in \mathcal{I}} \max_{\substack{S_1 \in \text{MIOS} \\ S_2 \in \text{DIS}}} \sum_k P(k \rightarrow r \mid \text{CFA}) p_k(S_1[\Theta], S_2[\Lambda]; \sigma, \mathbb{P})$$

Lower Bound

Theorem

The success probability of the order-finding protocol with operations Θ and unital Λ for creation and detection respectively is bounded by

$$P^{\text{succ}}(\Theta, \Lambda) \geq \frac{4}{\pi^2} \left(\frac{\varphi(r)}{r} \right) \left[\frac{1 + \mathcal{C}(\Theta) \tilde{M}_\diamond(\Lambda)}{2} \right]^L$$

- Euler's totient function

$$\frac{\varphi(r)}{r} \approx \frac{\text{const.}}{\log \log r}$$

- Cohering power $\mathcal{C}(\Theta) = \max_{\sigma \in \mathcal{I}} \mathcal{C}(\Theta(\sigma))$

$$\mathcal{C}(\rho) = \min_{\tau \in \mathcal{D}} \left\{ s \geq 0 : \frac{\rho + s\tau}{1 + s} \in \mathcal{I} \right\}$$

- NSID measure

$$\tilde{M}_\diamond(\Lambda) = \min_{\Phi \in \mathcal{DI}} \max_{\sigma} \|\Delta(\Lambda - \Phi)\sigma\|_1$$

$$0 \leq \mathcal{C}(\Theta), \tilde{M}_\diamond(\Lambda) \leq 1$$

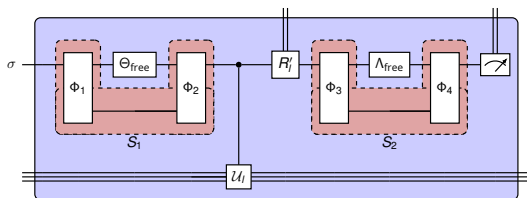
Theorem

The success probability of the order-finding protocol with operations Θ and unital Λ for creation and detection respectively is bounded by

$$P^{succ}(\Theta, \Lambda) \leq \varphi(r) \left(1 + 2 \left\lfloor \frac{2^L}{r^2} \right\rfloor \right) \left[\frac{1 + \mathcal{C}(\Theta) \tilde{M}_\diamond(\Lambda)}{2} \right]^L$$

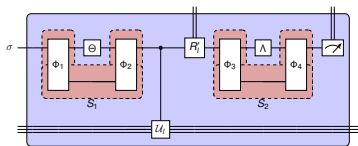
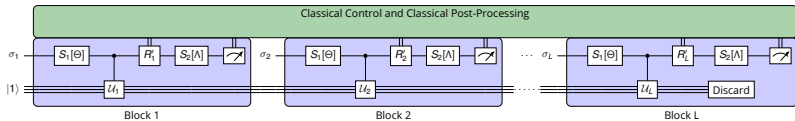
- Non-trivial in 'most' cases?
- Limitations of the prefactor?

How to define the classical limit?



$$\frac{2\varphi(r)}{2^L} \left\lfloor \frac{2^L - 1}{2r^2} \right\rfloor \leq P^{\text{succ}}(\Theta, \Lambda) \leq \frac{\varphi(r)}{2^L} \left(1 + 2 \left\lfloor \frac{2^L}{r^2} \right\rfloor \right)$$

Take-home message



$$P^{\text{succ}}(\Theta, \Lambda) \sim \left[\frac{1 + \mathcal{C}(\Theta) \tilde{M}_\diamond(\Lambda)}{2} \right]^L$$

Entanglement and coherence in Bernstein-Vazirani algorithm

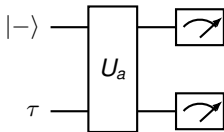
Moein Naseri, Tulja Varun Kondra, Suchetana Goswami, Marco Fellous-Asiani, and Alexander Streltsov
*Centre for Quantum Optical Technologies, Centre of New Technologies,
University of Warsaw, Banacha 2c, 02-097 Warsaw, Poland*

- Find bit string a

$$f(x) = a \cdot x = \left(\sum_{k=0}^N x_k a_k \right) \bmod 2$$

- Via an oracle

$$U_a |i\rangle |x\rangle = |i \oplus f(x)\rangle |x\rangle$$



- Optimal measurement

$$P^{\text{succ}}(\rho) = \max_{\{M_a\}} \sum_a \frac{1}{2^N} \text{Tr} [\mathcal{U}_a(\rho) M_a]$$

- NO entanglement required

$$P^{\text{succ}}(|-\rangle\langle-| \otimes \tau) = \frac{1 + C(\tau)}{2^N}$$

Need for a dynamical approach?