The contextual fraction and contextuality as a resource



Samson Abramsky¹



Rui Soares Barbosa¹



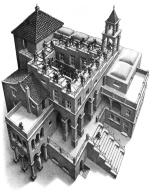
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Workshop on Quantum Contextuality in Quantum Mechanics and Beyond Prague, 4th June 2017

► Contextuality: a fundamental non-classical phenomenon of QM



- Contextuality: a fundamental non-classical phenomenon of QM
- Contextuality as a resource for QI and QC:
 - Non-local games quantum homomorphisms, constraint satisfaction, etc.
 - ▶ MBQC Raussendorf (2013) "Contextuality in measurement-based quantum computation"
 - MSD Howard, Wallman, Veith, & Emerson (2014) "Contextuality supplies the 'magic' for quantum computation"

► Abramsky—Brandenburger: unified framework for non-locality and contextuality in general measurement scenarios

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- qualitative hierarchy of contextuality for empirical models

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- qualitative hierarchy of contextuality for empirical models
- quantitative grading measure of contextuality (NB: there may be more than one useful measure)

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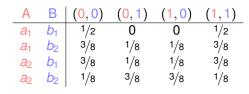
- ► **General**, i.e. applicable to any measurement scenario
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- Precise relationship to violations of Bell inequalities
- ► Monotone wrt operations that don't introduce contextuality

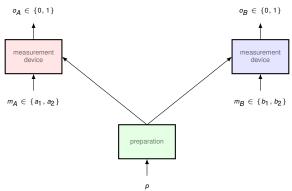
 → resource theory

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- Computable using linear programming
- Precise relationship to violations of Bell inequalities
- Relates to quantifiable advantages in QC and QIP tasks

Empirical data





Abramsky-Brandenburger framework

Measurement scenario $\langle X, \mathcal{M}, \mathcal{O} \rangle$:

- X is a finite set of measurements or variables
- O is a finite set of outcomes or values
- \triangleright \mathcal{M} is a cover of X, indicating **joint measurability** (contexts)

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Example: (2,2,2) Bell scenario

- ▶ The set of variables is $X = \{a_1, a_2, b_1, b_2\}$.
- ▶ The outcomes are $O = \{0, 1\}$.
- The measurement contexts are:

$$\{ \{a_1, b_1\}, \{a_1, b_2\}, \{a_2, b_1\}, \{a_2, b_2\} \}.$$

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A joint outcome or **event** in a context C is $s \in O^C$, e.g.

$$s = [a_1 \mapsto 0, b_1 \mapsto 1]$$
.

(These correspond to the cells of our probability tables.)

Another example: 18-vector Kochen-Specker

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- ▶ A set of 18 variables, X = {A,..., O}
- ▶ A set of outcomes *O* = {0, 1}
- A measurement cover $\mathcal{M} = \{C_1, \dots, C_9\}$, whose contexts C_i correspond to the columns in the following table:

U_1	U_2	U ₃	U_4	U_5	U_6	U_7	U ₈	U ₉
Α	Α	Н	Н	В	1	P	Р	Q
В	Ε	1	K	Ε	K	Q	R	R
С	F	С	G	М	Ν	D	F	М
D	G	J	L	Ν	0	J	L	0

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Compatibility condition: these distributions "agree on overlaps", i.e.

$$\forall_{C,C'\in\mathcal{M}}.\ e_C|_{C\cap C'} = e_{C'}|_{C\cap C'}.$$

where marginalisation of distributions: if $D \subseteq C$, $d \in Prob(O^C)$,

$$d|_{D}(s) := \sum_{t \in O^{C}, t|_{D} = s} d(t).$$

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For multipartite scenarios, compatibility = the **no-signalling** principle.

A (compatible) empirical model is **non-contextual** if there exists a **global distribution** $d \in \text{Prob}(O^X)$ (on the joint assignments of outcomes to all measurements) that marginalises to all the e_C :

$$\exists_{d \in \mathsf{Prob}(O^X)}. \ \forall_{C \in \mathcal{M}}. \ d|_C = e_C .$$

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family of data which is locally consistent but globally inconsistent.

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The import of results such as Bell's and Bell–Kochen–Specker's theorems is that there are empirical models arising from quantum mechanics that are contextual.

Strong contextuality

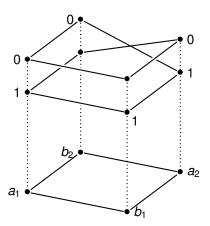
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E.g. K-S, GHZ, the PR box:

Α	В	(0,0)	(0, 1)	(1,0)	(1,1)
a ₁	b_1	✓	×	×	✓
		\checkmark	×	×	\checkmark
a_2	b_1	\checkmark	×	×	\checkmark
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Equivalently, maximum weight λ over all convex decompositions

$$e = \lambda e^{NC} + (1 - \lambda)e'$$

where e^{NC} is a non-contextual model.

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$$NCF(e) = \lambda$$
 $CF(e) = 1 - \lambda$

(Non-)contextual fraction via linear programming

Checking contextuality of e corresponds to solving

```
Find \mathbf{d} \in \mathbb{R}^n such that \mathbf{M} \, \mathbf{d} = \mathbf{v}^e and \mathbf{d} \geq \mathbf{0} .
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Computing the non-contextual fraction corresponds to solving the following linear program:

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Find \mathbf{c} \in \mathbb{R}^n maximising \mathbf{1} \cdot \mathbf{c} subject to \mathbf{M} \mathbf{c} \leq \mathbf{v}^e and \mathbf{c} \geq \mathbf{0} .
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E.g. Equatorial measurements on GHZ(n)

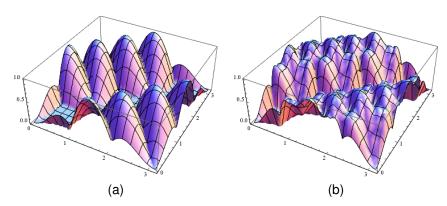


Figure: Contextual fraction of empirical models obtained with equatorial measurements at ϕ_1 and ϕ_2 on each qubit of $|\psi_{\text{GHZ}(n)}\rangle$ with: (a) n=3; (b) n=4.

Violations of Bell inequalities

An **inequality** for a scenario $\langle X, \mathcal{M}, \mathcal{O} \rangle$ is given by:

- ▶ a set of coefficients $\alpha = \{\alpha(C, s)\}_{C \in \mathcal{M}. s \in O^C}$
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For a model e, the inequality reads as

$$\mathcal{B}_{\alpha}(e) \leq R$$
,

where

$$\mathcal{B}_{\alpha}(\mathbf{e}) \; := \; \sum_{\mathbf{C} \in \mathcal{M}, \mathbf{s} \in \mathcal{O}^{\mathcal{C}}} \; \alpha(\mathbf{C}, \mathbf{s}) \mathbf{e}_{\mathbf{C}}(\mathbf{s}) \; .$$

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It is called a **Bell inequality** if it is satisfied by every NC model. If it is saturated by some NC model, the Bell inequality is said to be **tight**.

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A Bell inequality establishes a bound for the value of $\mathcal{B}_{\alpha}(e)$ amongst NC models.

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The **normalised violation** of a Bell inequality $\langle \alpha, R \rangle$ by an empirical model e is the value

$$\frac{\max\{0,\mathcal{B}_{\alpha}(e)-R\}}{\|\alpha\|-R}\;.$$

Proposition

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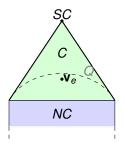
- ► The normalised violation by *e* of any Bell inequality is at most CF(*e*).
- ► This bound is attained: there exists a Bell inequality whose normalised violation by e is exactly CF(e).
- Moreover, this Bell inequality is tight at "the" non-contextual model e^{NC} and maximally violated by "the" strongly contextual model e^{SC} for any decomposition:

$$e = NCF(e)e^{NC} + CF(e)e^{SC}$$
.

Quantifying Contextuality LP:

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Find \mathbf{c} \in \mathbb{R}^n maximising \mathbf{1} \cdot \mathbf{c} subject to \mathbf{M} \, \mathbf{c} \leq \mathbf{v}^e and \mathbf{c} \geq \mathbf{0} .
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and

 $c\,\geq\,0$

Dual LP:

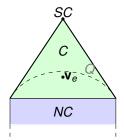
Find $\mathbf{y} \in \mathbb{R}^m$

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subject to $\mathbf{M}^T\mathbf{y} \geq \mathbf{1}$

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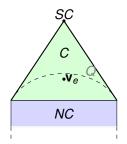
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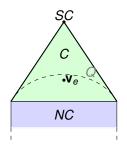
maximising $1 \cdot c$

subject to $Mc \leq v^e$

and

c > 0

$$e = \lambda e^{NC} + (1 - \lambda)e^{SC}$$
 with $\lambda = \mathbf{1} \cdot \mathbf{x}^*$.



Dual LP:

Find

 $\mathbf{v} \in \mathbb{R}^m$

minimising $\mathbf{v} \cdot \mathbf{v}^e$

subject to $\mathbf{M}^T \mathbf{y} \geq \mathbf{1}$

and

 $\mathbf{v} > \mathbf{0}$

$$\mathbf{a} := \mathbf{1} - |\mathcal{M}|\mathbf{y}$$

Find

 $\mathbf{a} \in \mathbb{R}^m$

maximising $\mathbf{a} \cdot \mathbf{v}^e$

subject to $\mathbf{M}^T \mathbf{a} < \mathbf{0}$

and

a < 1

Quantifying Contextuality LP:

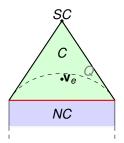
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Dual LP:

Find $y \in \mathbb{R}^m$ minimising $y \cdot v^e$ subject to $M^T y \ge 1$ and y > 0

 $a:=1-|\mathcal{M}|y$

Find $\mathbf{a} \in \mathbb{R}^m$ maximising $\mathbf{a} \cdot \mathbf{v}^e$ subject to $\mathbf{M}^T \mathbf{a} \leq \mathbf{0}$ and $\mathbf{a} < \mathbf{1}$.

computes tight Bell inequality (separating hyperplane)

Operations on empirical models

More than one possible measure of contextuality.

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- Towards a resource theory as for entanglement (e.g. LOCC), non-locality, ...

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The operations remind one of process algebras.

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Relabelling $\begin{array}{c} \mathbf{e}: \langle X, \mathcal{M}, \mathbf{O} \rangle \\ \alpha: (X, \mathcal{M}) \cong (X', \mathbf{M}') \end{array} \rightsquigarrow \mathbf{e}[\alpha]: \langle X', \mathcal{M}', \mathbf{O} \rangle$

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For
$$C \in \mathcal{M}, s : \alpha(C) \longrightarrow O$$
, $e[\alpha]_{\alpha(C)}(s) := e_C(s \circ \alpha^{-1})$

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Restriction

$$\begin{array}{ll} e: \langle X, \mathcal{M}, \mathcal{O} \rangle \\ (X', \mathcal{M}') \leq (X, M) & \rightsquigarrow & e \upharpoonright \mathcal{M}': \langle X', \mathcal{M}', \mathcal{O} \rangle \end{array}$$

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$$\boxed{ \mathsf{For} \; \pmb{C} \in \mathcal{M}, \pmb{s} : \alpha(\pmb{C}) \longrightarrow \pmb{O}, \, \pmb{e}[\alpha]_{\alpha(\pmb{C})}(\pmb{s}) := \pmb{e}_{\pmb{C}}(\pmb{s} \circ \alpha^{-1}) }$$

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$$\begin{array}{ll} e: \langle X, \mathcal{M}, \mathcal{O} \rangle \\ (X', \mathcal{M}') \leq (X, M) \end{array} \ \rightsquigarrow \ e \upharpoonright \mathcal{M}': \langle X', \mathcal{M}', \mathcal{O} \rangle \\ \end{array}$$

For
$$C' \in M', s : C' \longrightarrow O$$
, $(e \upharpoonright \mathcal{M}')_{C'}(s) := e_C|_{C'}(s)$ with any $C \in \mathcal{M}$ s.t. $C' \subseteq C$

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Coarse-graining

$$e: \langle X, \mathcal{M}, O \rangle$$

 $f: O \longrightarrow O'$ $\leadsto e/f: \langle X, \mathcal{M}, O' \rangle$

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For
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, $(e/f)_C(s) := \sum_{t:C \longrightarrow O, f \circ t = s} e_C(t)$

$$\begin{array}{ll} \textbf{e}, \textbf{e}' : \langle \textbf{X}, \mathcal{M}, \textbf{O} \rangle \\ \lambda \in [0,1] \end{array} \quad \rightsquigarrow \quad \textbf{e} +_{\lambda} \textbf{e}' : \langle \textbf{X}, \mathcal{M}, \textbf{O} \rangle \label{eq:epsilon}$$

Mixing

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$$\begin{array}{ll} e: \langle X, \mathcal{M}, O \rangle \\ e': \langle X', \mathcal{M}', O \rangle \end{array} \ \, \rightsquigarrow \ \, e \,\&\, \, e': \langle X \sqcup X', \mathcal{M} \sqcup \mathcal{M}', O \rangle \end{array}$$

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$$\begin{split} \mathcal{M}\star\mathcal{M}' &:= \{\textit{C} \sqcup \textit{D} \mid \textit{C} \in \mathcal{M}, \textit{D} \in \mathcal{M}'\} \\ \text{For } \textit{C} \in \mathcal{M}, \textit{D} \in \mathcal{M}', \textit{s} &= \langle \textit{s}_1, \textit{s}_2 \rangle : \textit{C} \sqcup \textit{D} \longrightarrow \textit{O}, \\ &(\textit{e} \otimes \textit{e}')_{\textit{C} \sqcup \textit{D}} \langle \textit{s}_1, \textit{s}_2 \rangle := \textit{e}_{\textit{C}}(\textit{s}_1) \, \textit{e}'_{\textit{D}}(\textit{s}_2) \end{split}$$

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Contextual fraction and quantum advantages

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- ▶ Measure of contextuality → to quantify such advantages.

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- ▶ Then, $1 \bar{p}_S \ge NCF(e)\nu(f)$.

Constraint system $\langle V, D, \Gamma \rangle$

- V finite set of variables
- D finite domain of values
- Γ finite set of formulae on the variables in V

Write $V(\phi)$ for variables that occur in ϕ .

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 (cf. Abramsky–Hardy "Logical Bell inequalities")
- We have: $1 \bar{p}_S \leq \text{NCF} \frac{(n-k)}{n}$.

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- Connections with Contextuality-by-Default (Dzhafarov et al.)

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 - What (else) is this resource useful for?

Questions...

